

Automatic Generation of PMEs for Matrix Operations

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$$LU = A \quad \longrightarrow$$

Partition $A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$

where A_{TL} is 0×0

While $n(A_{TL}) < n(A)$ **do**

$$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

$$A_{01} = L_{00}^{-1} A_{01}$$

$$A_{10} = A_{10} U_{01}^{-1}$$

$$A_{11} = LU(A_{11} - A_{10} A_{01})$$

$$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

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- We aim at loop-based algorithms

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- We aim at loop-based algorithms
- Correct by construction

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endwhile

- We aim at loop-based algorithms
- Loop invariants needed beforehand
- Correct by construction

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endwhile

- We aim at loop-based algorithms
- Correct by construction
- Loop invariants needed beforehand
- Loop invariants come from the PME



Cholesky Factorization

$$LL^T = A$$



Cholesky Factorization

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or in explicit form:

$$L := \Gamma(A)$$



Cholesky Factorization

$$LL^T = A$$

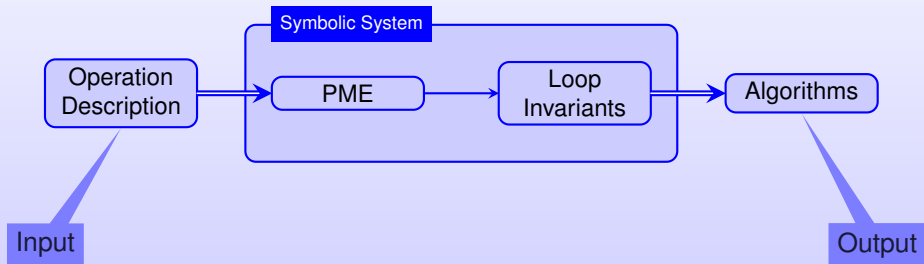
or in explicit form:

$$L := \Gamma(A)$$

Partitioned Matrix Expression:

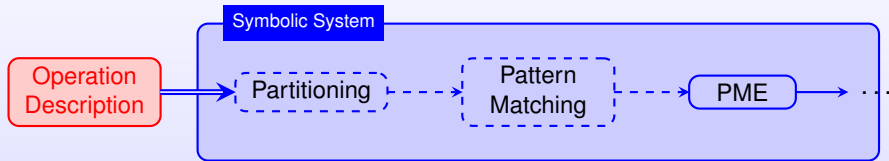
$$\left(\begin{array}{c|c} L_{TL} = \Gamma(A_{TL}) & * \\ \hline L_{BL} = A_{BL}L_{TL}^{-T} & L_{BR} = \Gamma(A_{BR} - L_{BL}L_{BL}^T) \end{array} \right)$$

Mechanical process: can be automated

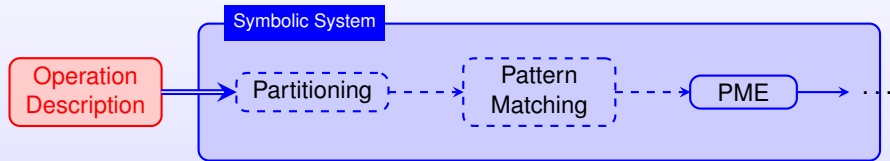


- 1 Motivation
- 2 Automatic Generation of PME**s
- 3 PME Examples
- 4 Conclusions

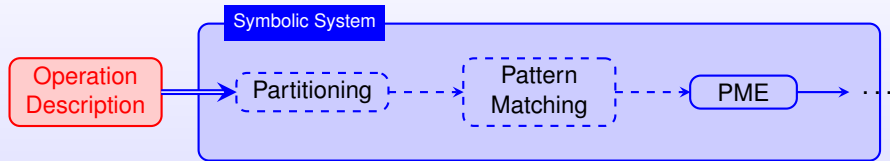




$$LU = A ?$$



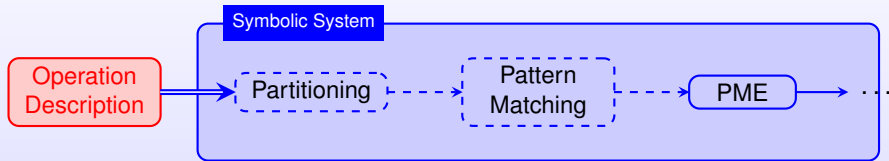
$$XY = Z ?$$



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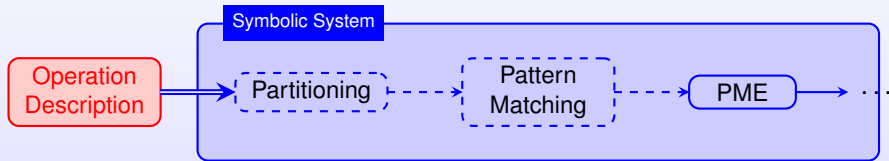
$$Z \leftarrow XY$$





$$XY = Z ?$$

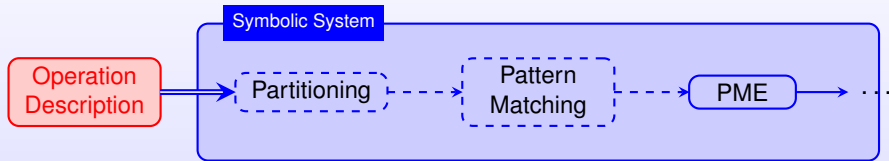
The diagram shows the equation $XY = Z ?$ at the top. Three arrows point downwards from a central point below the question mark to three different equations: $Z \leftarrow XY$, $X \leftarrow ZY^{-1}$, and $Y \leftarrow X^{-1}Z$.



$$XY = Z ?$$

Four arrows point downwards from the equation above to the following expressions:

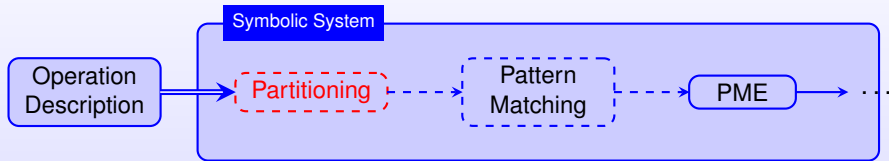
$$Z \leftarrow XY \quad X \leftarrow ZY^{-1} \quad Y \leftarrow X^{-1}Z \quad XY \leftarrow Z$$



$$LU = A$$

⇓

$$f : \{L, U\} := LU(A) \equiv \begin{cases} f_{\text{Pre}} : \{ \text{Input}(A) \wedge \text{Exists}LU(A) \wedge \\ \text{Output}(L) \wedge \text{LowTriUnit}(L) \wedge \\ \text{Output}(U) \wedge \text{UppTri}(U) \} \\ f_{\text{Post}} : \{LU = A\} \end{cases}$$

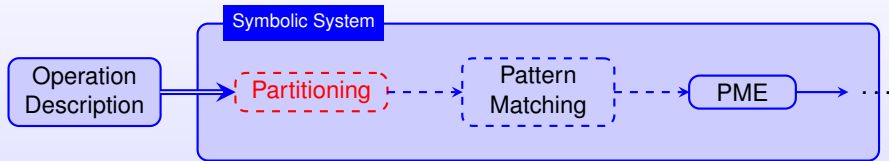


$$L \rightarrow (L)$$

$$L \rightarrow (L_L \mid L_R)$$

$$L \rightarrow \left(\frac{L_T}{L_B} \right)$$

$$L \rightarrow \left(\begin{array}{c|c} L_{TL} & L_{TR} \\ \hline L_{BL} & L_{BR} \end{array} \right)$$

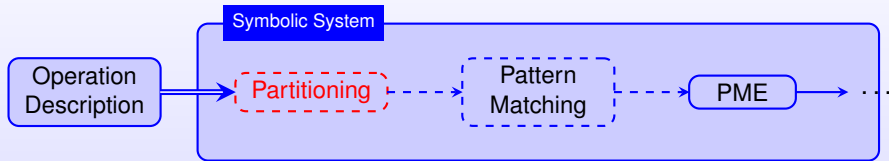


$L \rightarrow (L)$ where L is lower triangular

$L \rightarrow \left(\begin{array}{c|c} L_L & L_R \end{array} \right)$

$L \rightarrow \left(\begin{array}{c} L_T \\ \hline L_B \end{array} \right)$

$L \rightarrow \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right)$ where L_{TL} & L_{BR} are lower triangular



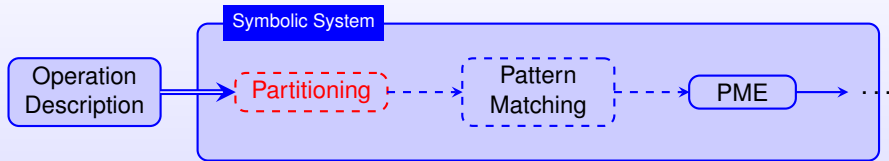
$L \rightarrow (L)$ where L is lower triangular

$L \rightarrow \left(\begin{array}{c|c} L_L & L_R \end{array} \right)$

$L \rightarrow \left(\begin{array}{c} L_T \\ \hline L_B \end{array} \right)$

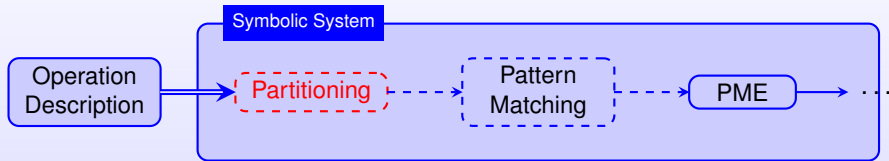
$L \rightarrow \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right)$ where L_{TL} & L_{BR} are lower triangular

Our system keeps track of the properties



Partitioned postcondition:

$$LU = A$$

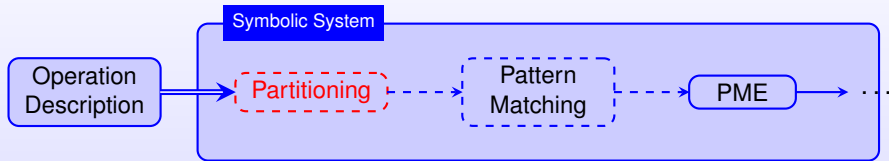


Partitioned postcondition:

$$LU = A$$

↓

$$\left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) (U) = \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$$



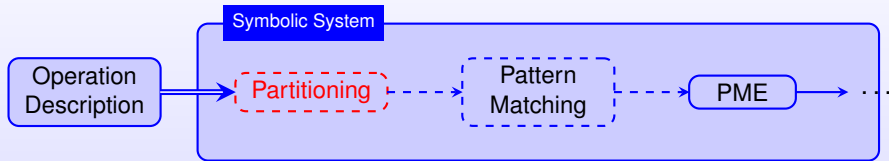
Partitioned postcondition:

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↓

$$\left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) (U) \equiv \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

Non-Conformal Partitioning

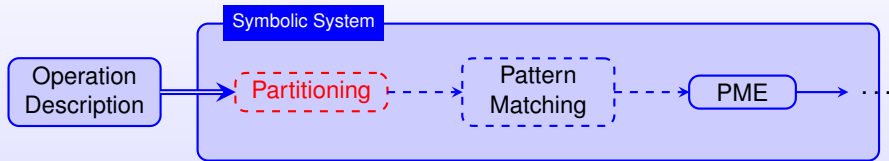


Partitioned postcondition:

$$LU = A$$

⇓

$$(L)(U) = (A)$$



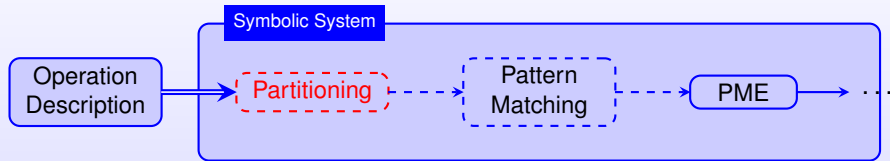
Partitioned postcondition:

$$LU = A$$

↓

$$\cancel{(L)(U) = (A)}$$

Does not decompose the operation

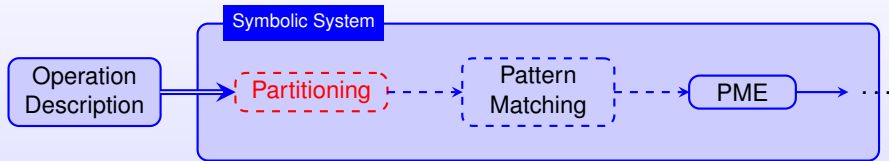


Partitioned postcondition:

$$LU = A$$

↓

$$\left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} U_{TL} & U_{TR} \\ \hline 0 & U_{BR} \end{array} \right) = \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$$



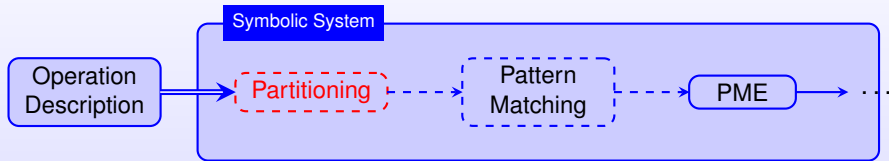
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Only the feasible ones are generated!

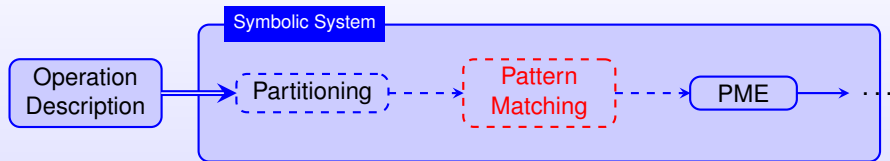


Symbolic computation:

$$\left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} U_{TL} & U_{TR} \\ \hline 0 & U_{BR} \end{array} \right) = \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

⇓

$$\left(\begin{array}{c|c} L_{TL}U_{TL} = A_{TL} & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$

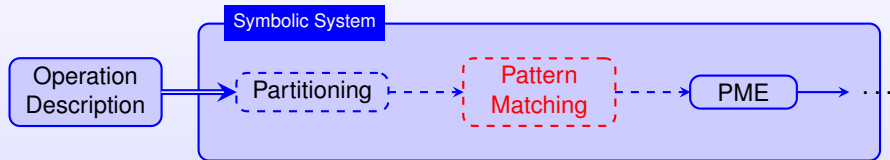


Input / Output:

$$\left(\begin{array}{c|c} L_{TL}U_{TL} = A_{TL} & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$

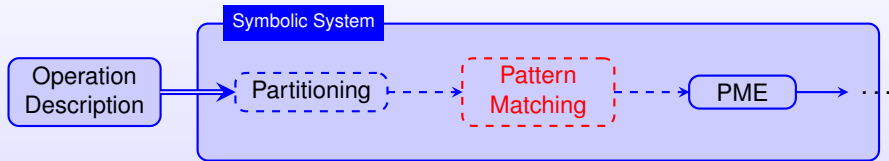
⇓

$$\left(\begin{array}{c|c} L_{TL}U_{TL} = A_{TL} & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$



Pattern matching and algebraic manipulation:

$$\left(\begin{array}{c|c} L_{TL}U_{TL} = A_{TL} & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right) \\ \Downarrow \\ \left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$

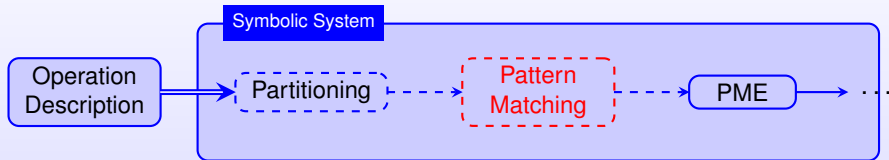


Pattern matching and algebraic manipulation:

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$

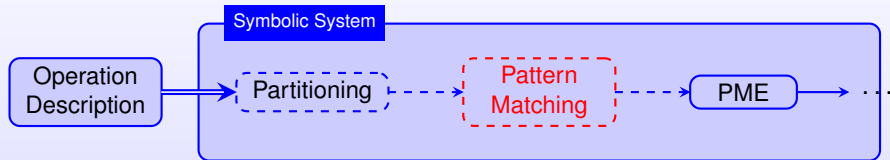
$$\Downarrow$$

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & L_{TL}U_{TR} = A_{TR} \\ \hline L_{BL}U_{TL} = A_{BL} & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$



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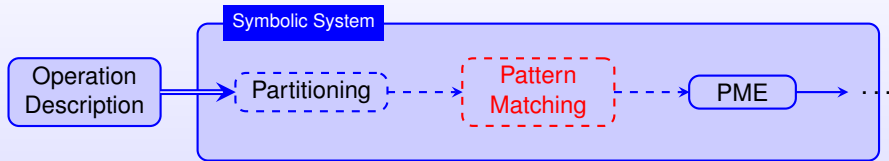


Pattern matching and algebraic manipulation:

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↓

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & U_{TR} = TRSM(L_{TL}^{-1}, A_{TR}) \\ \hline L_{BL} = TRSM(A_{BL}, U_{TL}^{-1}) & L_{BL}U_{TR} + L_{BR}U_{BR} = A_{BR} \end{array} \right)$$

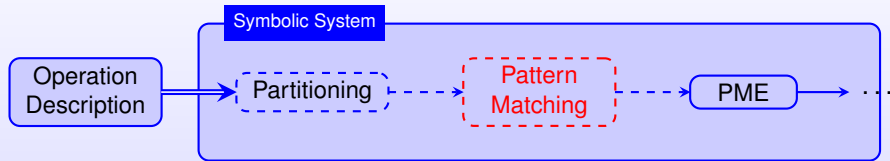


Pattern matching and algebraic manipulation:

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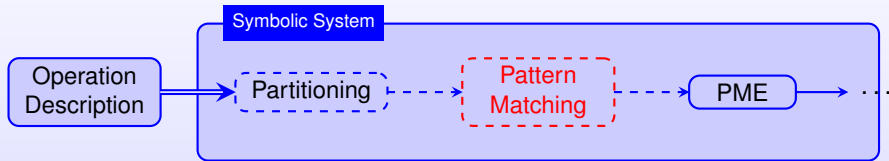


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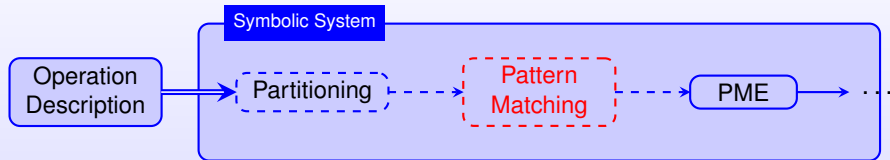
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Pattern matching and algebraic manipulation:

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & U_{TR} = TRSM(L_{TL}^{-1}, A_{TR}) \\ \hline L_{BL} = TRSM(A_{BL}, U_{TL}^{-1}) & L_{BR}U_{BR} = A_{BR} - L_{BL}U_{TR} \end{array} \right)$$



Pattern matching and algebraic manipulation:

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$ExistsLU(A_{BR} - L_{BL}U_{TR}) ?$

ExistsLU($A_{BR} - L_{BL}U_{TR}$) ?

Theorem

$$A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

$$\begin{array}{l} \text{ExistsLU}(A) \\ \text{full rank}(A) \end{array} \implies \left\{ \begin{array}{l} \text{ExistsLU}(A_{TL}) \wedge \\ \text{ExistsLU}(A_{BR} - A_{BL}A_{TL}^{-1}A_{TR}) \end{array} \right.$$

ExistsLU($A_{BR} - L_{BL}U_{TR}$) ?

Theorem

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ExistsLU($A_{BR} - L_{BL}U_{TR}$) \equiv ExistsLU($A_{BR} - A_{BL}A_{TL}^{-1}A_{TR}$) ?

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Theorem

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$$\begin{array}{l} \text{ExistsLU}(A) \\ \text{full rank}(A) \end{array} \implies \left\{ \begin{array}{l} \text{ExistsLU}(A_{TL}) \wedge \\ \text{ExistsLU}(A_{BR} - A_{BL}A_{TL}^{-1}A_{TR}) \end{array} \right.$$

ExistsLU($A_{BR} - L_{BL}U_{TR}$) \equiv ExistsLU($A_{BR} - A_{BL}A_{TL}^{-1}A_{TR}$) ?

$$L_{BL} \rightarrow A_{BL}U_{TL}^{-1} ; U_{TR} \rightarrow L_{TL}^{-1}A_{TR} ; L_{TL} * U_{TL} \rightarrow A_{TL}$$

ExistsLU($A_{BR} - L_{BL}U_{TR}$) ?

Theorem

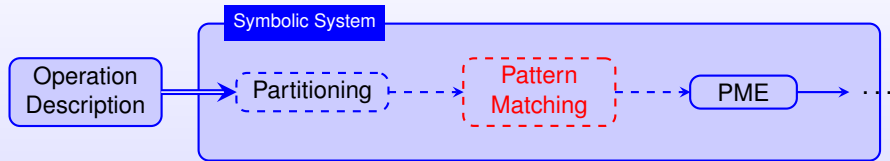
$$A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

$$\begin{array}{l} \text{ExistsLU}(A) \\ \text{full rank}(A) \end{array} \implies \left\{ \begin{array}{l} \text{ExistsLU}(A_{TL}) \wedge \\ \text{ExistsLU}(A_{BR} - A_{BL}A_{TL}^{-1}A_{TR}) \end{array} \right.$$

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$$L_{BL} \rightarrow A_{BL}U_{TL}^{-1} ; U_{TR} \rightarrow L_{TL}^{-1}A_{TR} ; L_{TL} * U_{TL} \rightarrow A_{TL}$$

$$L_{BL}U_{TR} \equiv A_{BL}U_{TL}^{-1}L_{TL}^{-1}A_{TR} \equiv A_{BL}(L_{TL}U_{TL})^{-1}A_{TR} \equiv A_{BL}A_{TL}^{-1}A_{TR}$$

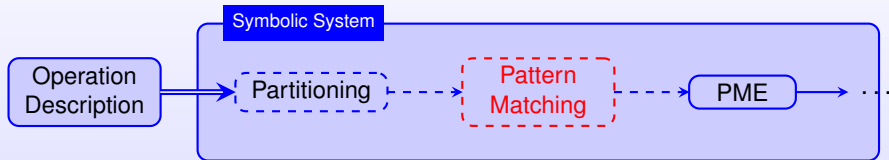


Pattern matching and algebraic manipulation:

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & U_{TR} = TRSM(L_{TL}^{-1}, A_{TR}) \\ \hline L_{BL} = TRSM(A_{BL}, U_{TL}^{-1}) & L_{BR}U_{BR} = A_{BR} - L_{BL}U_{TR} \end{array} \right)$$

⇓

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & U_{TR} = TRSM(L_{TL}^{-1}, A_{TR}) \\ \hline L_{BL} = TRSM(A_{BL}, U_{TL}^{-1}) & \{L_{BR}, U_{BR}\} = LU(A_{BR} - L_{BL}U_{TR}) \end{array} \right)$$



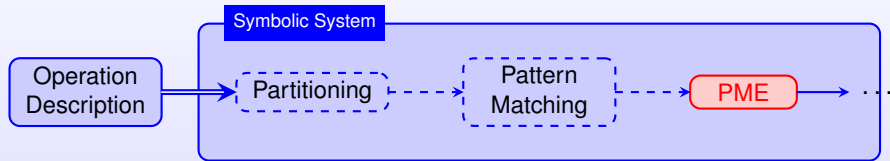
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PME:

$$\left(\begin{array}{c|c} \{L_{TL}, U_{TL}\} = LU(A_{TL}) & U_{TR} = TRSM(L_{TL}^{-1}, A_{TR}) \\ \hline L_{BL} = TRSM(A_{BL}, U_{TL}^{-1}) & \{L_{BR}, U_{BR}\} = LU(A_{BR} - L_{BL}U_{TR}) \end{array} \right)$$

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Sylvester Equation:

$$LX + XU = C$$

$$f : X := \Omega(L, U, C) \equiv \begin{cases} f_{\text{Pre}} : \{ \text{Input}(L, U, C) \wedge \text{Output}(X) \wedge \\ \text{LowTri}(L) \wedge \text{UppTri}(U) \} \\ f_{\text{Post}} : \{ LX + XU = C \} \end{cases}$$

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↓

$$\begin{pmatrix} C_T \\ C_B \end{pmatrix} = \begin{pmatrix} L_{TL} & | & 0 \\ L_{BL} & | & L_{BR} \end{pmatrix} \begin{pmatrix} X_T \\ X_B \end{pmatrix} + \begin{pmatrix} X_T \\ X_B \end{pmatrix} U$$

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$$\begin{pmatrix} X_T = \Omega(L_{TL}, U, C_T) \\ X_B = \Omega(L_{BR}, U, C_B - L_{BL}X_T) \end{pmatrix}$$

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$$(C_L \mid C_R) = L (X_L \mid X_R) + (X_L \mid X_R) \left(\begin{array}{c|c} U_{TL} & U_{TR} \\ \hline 0 & U_{BR} \end{array} \right)$$

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⇓

$$(X_L = \Omega(L, U_{TL}, C_L) \mid X_R = \Omega(L, U_{BR}, C_R - X_L U_{TR}))$$

Sylvester Equation:

$$LX + XU = C$$

$$f : X := \Omega(L, U, C) \equiv \begin{cases} f_{\text{Pre}} : \{ \text{Input}(L, U, C) \wedge \text{Output}(X) \wedge \\ \text{LowTri}(L) \wedge \text{UppTri}(U) \} \\ f_{\text{Post}} : \{ LX + XU = C \} \end{cases}$$

↓

$$\left(\begin{array}{c|c} C_{TL} & C_{TR} \\ \hline C_{BL} & C_{BR} \end{array} \right) = \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} X_{TL} & X_{TR} \\ \hline X_{BL} & X_{BR} \end{array} \right) + \left(\begin{array}{c|c} X_{TL} & X_{TR} \\ \hline X_{BL} & X_{BR} \end{array} \right) \left(\begin{array}{c|c} U_{TL} & U_{TR} \\ \hline 0 & U_{BR} \end{array} \right)$$

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⇓

$$\left(\begin{array}{c|c} X_{TL} = \Omega(L_{TL}, U_{TL}, C_{TL}) & X_{TR} = \Omega(L_{TL}, U_{BR}, C_{TR} - X_{TL}U_{TR}) \\ \hline X_{BL} = \Omega(L_{BR}, U_{TL}, C_{BL} - L_{BL}X_{TL}) & X_{BR} = \Omega(L_{BR}, U_{BR}, C_{BR} - L_{BL}X_{TR} - X_{BL}U_{TR}) \end{array} \right)$$

Reduction to Standard Form:

$$LBL^T = A$$

$$f : B := \Phi(A, L) \equiv \begin{cases} f_{\text{Pre}} : \{ \text{Input}(A, L) \wedge \text{Output}(B) \wedge \\ \quad \text{LowTri}(L) \wedge \text{Symmetric}(A, B) \} \\ f_{\text{Post}} : \{ LBL^T = A \} \end{cases}$$

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⇓

$$\left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} B_{TL} & B_{BL}^T \\ \hline B_{BL} & B_{BR} \end{array} \right) \left(\begin{array}{c|c} L_{TL}^T & L_{BL}^T \\ \hline 0 & L_{BR}^T \end{array} \right) = \left(\begin{array}{c|c} A_{TL} & A_{BL}^T \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

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$$\left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} B_{TL} & B_{BL}^T \\ \hline B_{BL} & B_{BR} \end{array} \right) \left(\begin{array}{c|c} L_{TL}^T & L_{BL}^T \\ \hline 0 & L_{BR}^T \end{array} \right) = \left(\begin{array}{c|c} A_{TL} & A_{BL}^T \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

⇓

$$\left(\begin{array}{c|c} B_{TL} = \Phi(A_{TL}, L_{TL}) & * \\ \hline B_{BL} = L_{BR}^{-1}(A_{BL}L_{TL}^{-1} - L_{BL}B_{TL}) & B_{BR} = \Phi(A_{BR} - L_{BL}B_{TL}L_{BL}^T - L_{BR}B_{BL}L_{BL}^T \\ & - L_{BL}B_{BL}^T L_{BR}^T, L_{BR}) \end{array} \right)$$

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- Minimum amount of knowledge:
 - Input and output operands
 - Size of the operands
 - Structure of the operands: triangularity, symmetry, ...



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- Knowledge implemented **IN** the system:
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 - Theorems
- One step closer to the full automatic generation of algorithms.



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