

Knowledge-Based Automatic Generation of Algorithms and Code

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Focus

Design and implementation of Domain-Specific Compilers for Linear Algebra matrix equations.

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Why?

- Matrix equations are ubiquitous

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- It requires expertise from multiple areas:
 - Application domain
 - Numerics, algorithmics
 - High-performance computing

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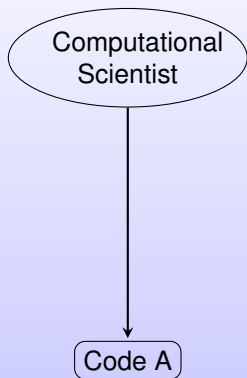
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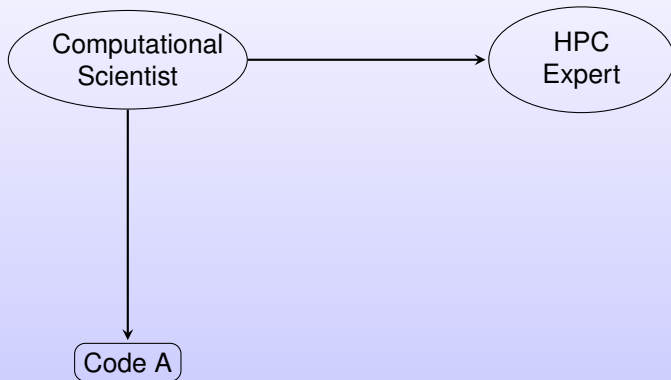
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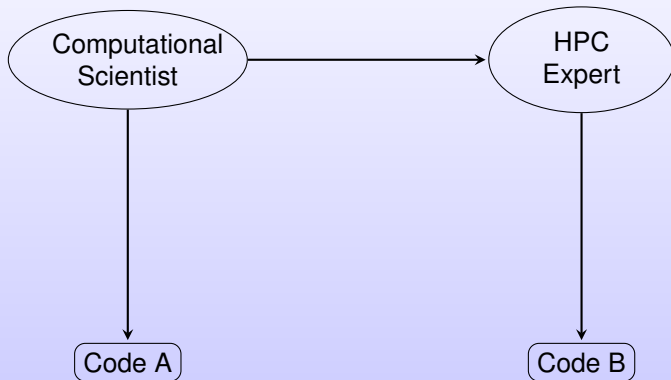
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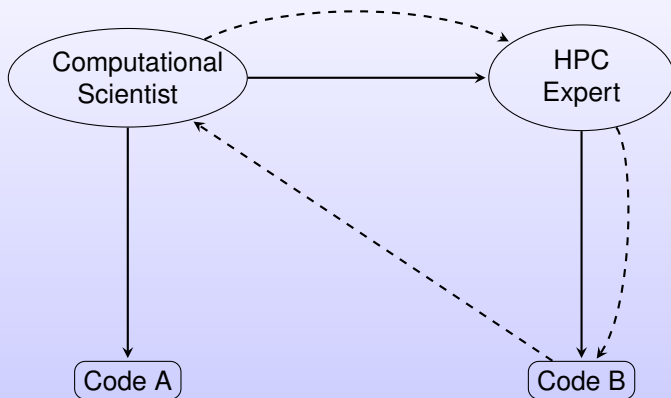
We are facing a productivity problem

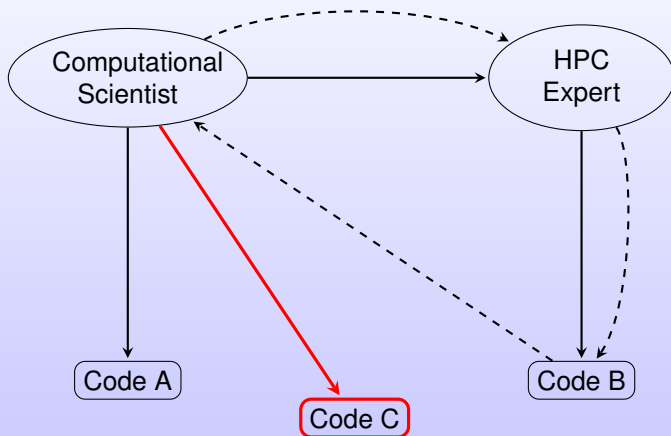
Computational
Scientist









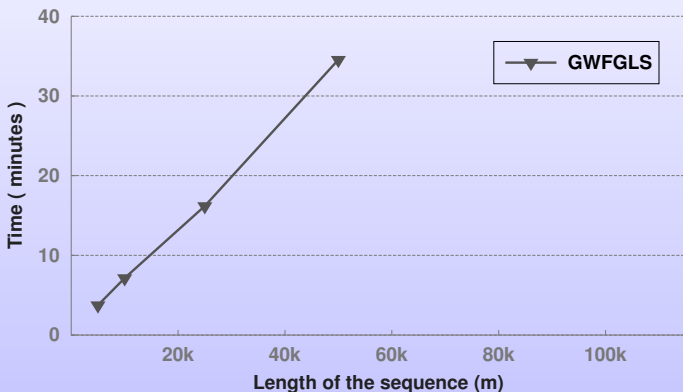


Example 1

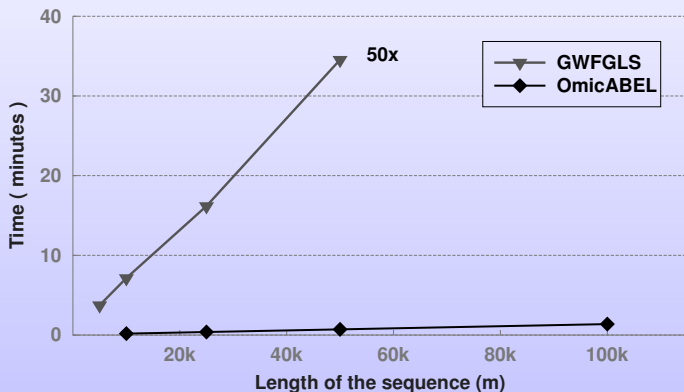
Genome-Wide Association Study

$$b_i := (X_i^T M^{-1} X_i)^{-1} X_i^T M^{-1} y \quad 1 \leq i \leq m$$
$$M \in \mathbb{R}^{n \times n}, X \in \mathbb{R}^{n \times p}, y \in \mathbb{R}^n, b \in \mathbb{R}^p$$

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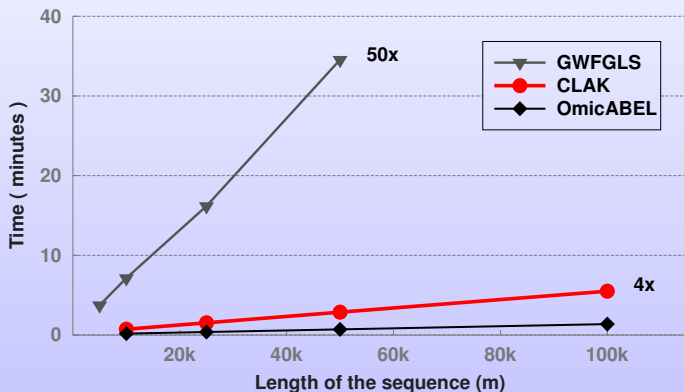


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Derivative of the Cholesky factorization

- $f : LL^T = A \rightarrow (\text{LAPACK, FLAME, ...})$

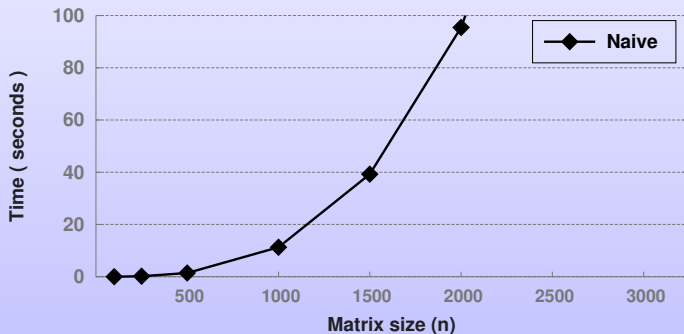
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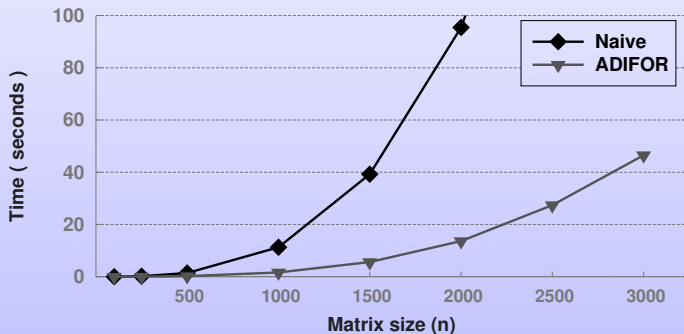
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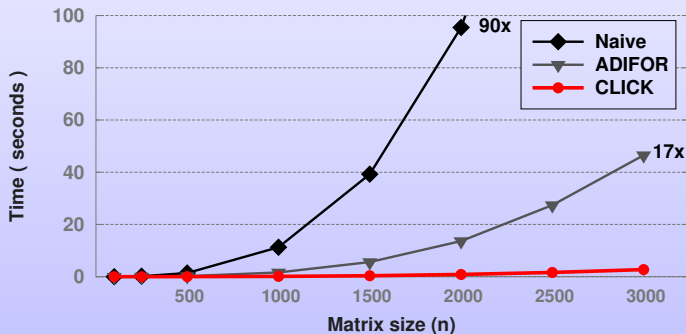
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Productivity (+ Performance)

1 Two Linear Algebra Compilers

2 CLAK

3 CL1CK

4 Contributions

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- Main idea: Decomposition onto building blocks

$$X := S^{-1} \quad \longrightarrow$$

Algorithm 2 : $X := S^{-1}$.

- 1: $LL^T = S$ (Cholesky factorization)
 - 2: $L := L^{-1}$ (Triangular inverse)
 - 3: $X := L^T L$ (Matrix product)
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- Methodology: Replicate the reasoning of a human expert

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- Core idea: Loop-based blocked algorithms

$$LL^T = A \quad \longrightarrow$$

$$\text{Partition } A \rightarrow \left(\begin{array}{c|c} A_{TL} & * \\ \hline A_{BL} & A_{BR} \end{array} \right)$$

where A_{TL} is 0×0

While $n(A_{TL}) < n(A)$ do

$$\left(\begin{array}{c|c} A_{TL} & * \\ \hline A_{BL} & A_{BR} \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} A_{00} & * & * \\ \hline A_{10} & A_{11} & * \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

$$A_{11} = \Gamma(A_{11})$$

$$A_{21} = A_{21} \text{TRIL}(A_{11})^{-T}$$

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- Core idea: Loop-based blocked algorithms
- Methodology: FLAME Project's methodology

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```
Equation SEQ_OLS
# Operands declaration
Vector b    <Output>;
Matrix X    <Input, FullRank, ColumnPanel>;
Vector y    <Input>;
# Equation
b{i} = inv( trans(X{i}) * X{i} ) * trans(X{i}) * y
```

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CLAK ...

- ... breaks a target operation down to **building blocks** ...
- ... directly supported by **high-performance libraries**, ...
- ... **tailoring** the algorithm **to the application**.

Traditional Comp.

- ADD
- MUL
- DIV
- XOR

Linear Algebra Comp.

- $LU = A$ (LU)
 - $LL^T = A$ (Cholesky)
 - $ZWZ^T = A$ (Eigendec)
-
- $C := AB$ (MM)
 - $y := Ax$ (MV)
 - $Ax = b$ (TRSV)

CLAK aims at replicating the reasoning of a human expert

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We ...

- studied the steps an expert takes,
- encoded them in a set of heuristics, and
- incorporated these heuristics into CLAK

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$$x := A^{-1}b$$

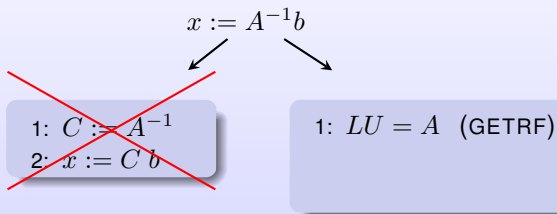
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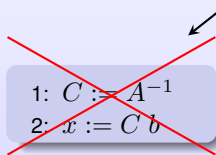
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2: $y := L^{-1}b$ (TRSV)
3: $x := U^{-1}y$ (TRSV)

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- Do not invert unless really required

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Computation reuse

- 1: $x := L^{-T} y$
- 2: $\alpha := x^T x$

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Algorithm 1

- 1: $C := AB$ $O(n^3)$
- 2: $y := Cx$ $O(n^2)$

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Algorithm 2

- 1: $t := Bx$ $O(n^2)$
- 2: $y := At$ $O(n^2)$

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Priorities: Matrix-Vector over Matrix-Matrix

CLAK applies these heuristics mechanically:

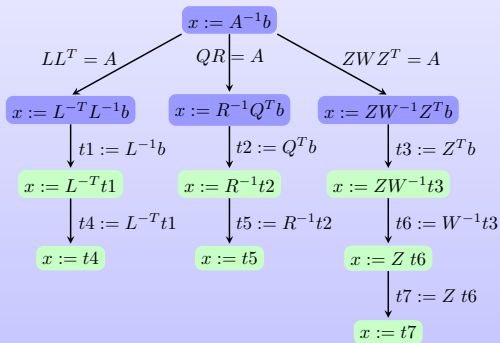
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As a result, CLAK generates a tree of decompositions:



Factorizations

Matrix Property	Factorizations
Symmetric	LDL, QR, Eigendecomposition
SPD	Cholesky, QR, Eigendecomposition
Column Panel	QR
...	...

Reducing complexity

- Prioritization based on dimensionality of the operands:

#	Kernels	Example
1	inner product	$\alpha := x^T y$
2	matrix-vector operations	$y := Ax, b := L^{-1}x$
3	matrix-matrix operations	$C := AB, B := L^{-1}A$
4	outer product	$A := xy^T$
5	inversion of a triangular matrix	$C := L^{-1}$

- Common segments ($\alpha := y^T L^{-1} L^{-T} y$)

Built in a modular fashion

- Matrix algebra
- Inference of knowledge
- Other modules

$$(A \times B)^T \rightarrow B^T \times A^T$$

$$(A \times B)^{-1} \wedge \text{Square}(A) \wedge \text{Square}(B) \rightarrow B^{-1} \times A^{-1}$$

$$Q^T \times Q \wedge \text{Orthogonal}(Q) \rightarrow I$$

$$A^{-1} \times A \rightarrow I$$

$$A \times I \wedge \text{Matrix}(A) \rightarrow A$$

$$\begin{aligned}(A \times B)^T &\rightarrow B^T \times A^T \\(A \times B)^{-1} \wedge \text{Square}(A) \wedge \text{Square}(B) &\rightarrow B^{-1} \times A^{-1} \\Q^T \times Q \wedge \text{Orthogonal}(Q) &\rightarrow I \\A^{-1} \times A &\rightarrow I \\A \times I \wedge \text{Matrix}(A) &\rightarrow A\end{aligned}$$

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$$\begin{aligned}((QR)^T QR)^{-1} (QR)^T L^{-1} y &\longrightarrow R^{-1} Q^T L^{-1} y \\(ZWZ^T + I)^{-1} &\longrightarrow Z(W + I)^{-1} Z^T\end{aligned}$$

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About 50 such rules

Type of operand:

Expression	Constraint	Inferred Property
$A \times B$	$\text{Matrix}(A) \wedge \text{Matrix}(B)$	$\text{Matrix}(A \times B)$
$A \times x$	$\text{Matrix}(A) \wedge \text{Vector}(x)$	$\text{Vector}(A \times x)$
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Matrix factorizations (properties of factors):

QR ($QR = A$):

Input A : matrix, column-panel, full rank

Output Q : matrix, orthogonal, column-panel, full rank

R : matrix, square, upper triangular, full rank

Building blocks:

Expression	Constraint	Inferred Property
$S_1 + \dots + S_n$	$\forall_i \text{Symmetric}(S_i)$	Symmetric
$-S$	$\text{Symmetric}(S)$	Symmetric
S^T	$\text{Symmetric}(S)$	Symmetric
S^{-1}	$\text{Symmetric}(S)$	Symmetric
expr	$\text{expr} == \text{expr}^T$	Symmetric

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More than a hundred such rules!

- Sequences of problems

Algorithm 4 CLAK-EIG

```
1:  $Z\Lambda Z^T = \Phi$ 
2: for  $i := 1$  to  $m$  do
3:    $K_i := X_i^T Z$ 
4: end for
5: for  $j := 1$  to  $t$  do
6:    $D_j := h_j \Lambda + (1 - h_j)I$ 
7:    $y_j := Z^T y_j$ 
8:   for  $i := 1$  to  $m$  do
9:      $V_{ij} := K_i D_j^{-1}$ 
10:     $A_{ij} := V_{ij} K_i^T$ 
11:     $Q_{ij} R_{ij} = A_{ij}$ 
12:     $b_{ij} := V_{ij} y_j$ 
13:     $b_{ij} := Q_{ij}^T b_{ij}$ 
14:     $b_{ij} := R_{ij}^{-1} b_{ij}$ 
15:   end for
16: end for
```

• Derivative Operator

$z = \text{alpha} * x + y \longrightarrow \text{dv}(z = \text{alpha} * x + y) ?:$

- $\text{dv}(z) = \text{dv}(\text{alpha}) * x + \text{alpha} * \text{dv}(x) + \text{dv}(y)$
- $\text{dv}(z) = \text{dv}(\text{alpha}) * x + \text{alpha} * \text{dv}(x)$
- $\text{dv}(z) = \text{dv}(\text{alpha}) * x$
- ...

• Matlab Code

```
1 function [b] = GWAS_26_2(X, y, h, Phi, sm, sn, nXs, nys)
2     b = zeros(sm, nXs * nys);
3     T3 = zeros(sm, sn * nXs);
4     [Z1, W1] = eig( Phi );
5
6     for i = 1:nXs
7         T3(:, sn*(i-1)+1:sn*i) = X(:, sm*(i-1)+1:sm*i)' * Z1;
8     end
9     for j = 1:nys
10        T1 = 1 * eye(sn) + - h(j) * eye(sn);
11        T2 = T1 + h(j) * W1;
12        T6 = Z1' * y(:, j);
13        for i = 1:nXs
14            [...]
```

• Fortran Code

```
1 SUBROUTINE GWAS_26_2( X, csX, dsX, y, csy, h, Phi, csPhi,
2                     b, csb, sn, sm, nXs, nys )
3   INTEGER sn, sm, nXs, nys, csX, dsX, csy, csPhi, csb
4   [...]
5   call dsyevr( 'V', 'A', 'L', sn, Phi( 1, 1 ), sn, ddummy,
6              ddummy, idummy, idummy, ddummy, nCompPairs5,
7              W8( 1 ), Z7( 1, 1 ), sn, isuppz4, ... )
8   DO i = 1, nXs
9     call dgemm( 'T', 'N', sm, sn, sn, ONE, X( 1, 1, i ),
10              sn, Z7( 1, 1 ), sn, ZERO, tmp70( 1, 1, i ), sm )
11   END DO
12   DO j = 1, nys
13     DO iter1 = 1, sn
14       tmp28( iter1 ) = 1 + ( - h(j))
15       [...]
```

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- Multiple linear systems
- Ordinary least-squares
- Generalized least-squares
- Sequences of problems

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- Equations arising in Genome-wide association studies
- Derivatives of matrix products
- Derivatives of linear systems

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FLAME Methodology

- More than a decade of development
- Many algorithms derived by hand and incorporated into libraries (FLAME, Elemental)
- Manual derivation becomes tedious and error prone when the complexity of the equations increases

CL1CK demonstrates...

- FLAME methodology can be applied automatically
- The broad applicability of the methodology

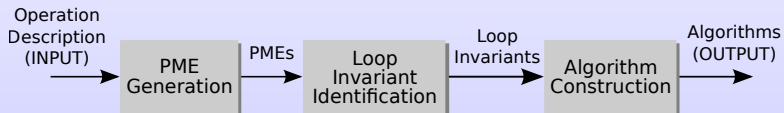
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Example: Derivative of Cholesky

$$G = gChol(L, B) \equiv \left\{ \begin{array}{l} P_{\text{pre}} : \{ \text{Output}(G) \wedge \text{Input}(L) \wedge \text{Input}(B) \wedge \\ \text{Matrix}(G) \wedge \text{Matrix}(L) \wedge \text{Matrix}(B) \wedge \\ \text{LowerTriangular}(G) \wedge \text{Symmetric}(B) \wedge \\ \text{LowerTriangular}(L) \} \\ P_{\text{post}} : \{ GL^T + LG^T = B \} \end{array} \right.$$

- CL1CK implements the FLAME methodology in 3 stages:



$$G = gChol(L, B) \equiv \left\{ \begin{array}{l} P_{\text{pre}} : \{\text{Output}(G) \wedge \text{Input}(L) \wedge \text{Input}(B) \wedge \\ \text{Matrix}(G) \wedge \text{Matrix}(L) \wedge \text{Matrix}(B) \wedge \\ \text{LowerTriangular}(G) \wedge \text{Symmetric}(B) \wedge \\ \text{LowerTriangular}(L)\} \\ P_{\text{post}} : \{GL^T + LG^T = B\} \end{array} \right.$$

$$G = gChol(L, B) \equiv \left\{ \begin{array}{l} P_{pre} : \{Output(G) \wedge Input(L) \wedge Input(B) \wedge \\ Matrix(G) \wedge Matrix(L) \wedge Matrix(B) \wedge \\ LowerTriangular(G) \wedge Symmetric(B) \wedge \\ LowerTriangular(L)\} \\ P_{post} : \{GL^T + LG^T = B\} \end{array} \right.$$

⇓

```
equal[
  plus[ times[ G_, trans[L_] ], times[ L_, trans[G_] ] ], B_
] /; isInputQ[L] && isInputQ[B] && isOutputQ[G] &&
  isMatrixQ[L] && isMatrixQ[B] && isMatrixQ[G] &&
  isLowerTriQ[L] && isSymmetricQ[B] && isLowerTriQ[G]
```

Partitioned Matrix Expressions are generated in 3 steps:

- 1 Decompose the problem into smaller ones
- 2 Find out how to solve the sub-problems
- 3 Combine the solutions

$$GL^T + LG^T = B$$

$$GL^T + LG^T = B$$

⇓

$$\left(\begin{array}{c|c} G_{TL} & 0 \\ \hline G_{BL} & G_{BR} \end{array} \right) \left(\begin{array}{c|c} L_{TL}^T & L_{BL}^T \\ \hline 0 & L_{BR}^T \end{array} \right) + \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} G_{TL}^T & G_{BL}^T \\ \hline 0 & G_{BR}^T \end{array} \right) = \left(\begin{array}{c|c} B_{TL} & B_{BL}^T \\ \hline B_{BL} & B_{BR} \end{array} \right)$$

$$GL^T + LG^T = B$$

$$\Downarrow$$

$$\left(\begin{array}{c|c} G_{TL} & 0 \\ \hline G_{BL} & G_{BR} \end{array} \right) \left(\begin{array}{c|c} L_{TL}^T & L_{BL}^T \\ \hline 0 & L_{BR}^T \end{array} \right) + \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \left(\begin{array}{c|c} G_{TL}^T & G_{BL}^T \\ \hline 0 & G_{BR}^T \end{array} \right) = \left(\begin{array}{c|c} B_{TL} & B_{BL}^T \\ \hline B_{BL} & B_{BR} \end{array} \right)$$

$$\Downarrow$$

$$\left(\begin{array}{c|c} G_{TL}L_{TL}^T + L_{TL}G_{TL}^T = B_{TL} & * \\ \hline G_{BL}L_{TL}^T + L_{BL}G_{TL}^T = B_{BL} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{l|l} G_{TL}L_{TL}^T + L_{TL}G_{TL}^T = B_{TL} & * \\ \hline G_{BL}L_{TL}^T + L_{BL}G_{TL}^T = B_{BL} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{l|l} G_{TL}L_{TL}^T + L_{TL}G_{TL}^T = B_{TL} & * \\ G_{BL}L_{TL}^T + L_{BL}G_{TL}^T = B_{BL} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

↓

$$\left(\begin{array}{l|l} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ G_{BL}L_{TL}^T + L_{BL}G_{TL}^T = B_{BL} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL}L_{TL}^T + L_{BL}G_{TL}^T = B_{BL} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL}L_{TL}^T + L_{BL}G_{TL}^T = B_{BL} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

↓

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL}L_{TL}^T = B_{BL} - L_{BL}G_{TL}^T & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL}L_{TL}^T = B_{BL} - L_{BL}G_{TL}^T & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL}L_{TL}^T = B_{BL} - L_{BL}G_{TL}^T & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

↓

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BL}L_{BL}^T + G_{BR}L_{BR}^T + L_{BL}G_{BL}^T + L_{BR}G_{BR}^T = B_{BR} \end{array} \right)$$

↓

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BR}L_{BR}^T + L_{BR}G_{BR}^T = B_{BR} - G_{BL}L_{BL}^T - L_{BL}G_{BL}^T \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BR}L_{BR}^T + L_{BR}G_{BR}^T = B_{BR} - G_{BL}L_{BL}^T - L_{BL}G_{BL}^T \end{array} \right)$$

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BR}L_{BR}^T + L_{BR}G_{BR}^T = B_{BR} - G_{BL}L_{BL}^T - L_{BL}G_{BL}^T \end{array} \right)$$

↓

$$\left(\begin{array}{c|c} G_{TL} := gChol(L_{TL}, B_{TL}) & * \\ \hline G_{BL} := (B_{BL} - L_{BL}G_{TL}^T)L_{TL}^{-T} & G_{BR} := gChol(L_{BR}, B_{BR} - G_{BL}L_{BL}^T - L_{BL}G_{BL}^T) \end{array} \right)$$

Triangular Sylvester Equation ($AX + XB = C$):

$$\left(\begin{array}{c|c} X_{TL} := \Omega(A_{TL}, B_{TL}, C_{TL}) & X_{TR} := \Omega(A_{TL}, B_{BR}, C_{TR} - X_{TL}B_{TR}) \\ \hline X_{BL} := \Omega(A_{BR}, B_{TL}, C_{BL} - A_{BL}X_{TL}) & X_{BR} := \Omega(A_{BR}, B_{BR}, \\ & C_{BR} - X_{BL}B_{TR} - A_{BL}X_{TR}) \end{array} \right)$$

Triangular Sylvester Equation ($AX + XB = C$):

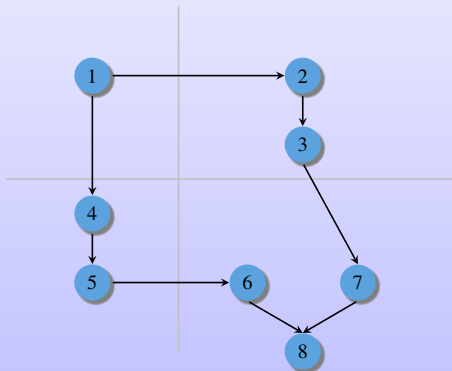
$$\left(\begin{array}{c|c} X_{TL} := \Omega(A_{TL}, B_{TL}, C_{TL}) & X_{TR} := \Omega(A_{TL}, B_{BR}, C_{TR} - X_{TL}B_{TR}) \\ \hline X_{BL} := \Omega(A_{BR}, B_{TL}, C_{BL} - A_{BL}X_{TL}) & X_{BR} := \Omega(A_{BR}, B_{BR}, \\ & C_{BR} - X_{BL}B_{TR} - A_{BL}X_{TR}) \end{array} \right)$$

Triangular Lyapunov Equation ($AX + XA^T = C$):

$$\left(\begin{array}{c|c} X_{TL} := \Lambda(A_{TL}, C_{TL}) & * \\ \hline X_{BL} := \Omega(A_{BR}, A_{TL}^T, C_{BL} - A_{BL}X_{TL}) & X_{BR} := \Lambda(A_{BR}, \\ & C_{BR} - X_{BL}A_{BL}^T - A_{BL}X_{BL}^T) \end{array} \right)$$

Graph of dependencies

$$\left(\begin{array}{c|c} X_{TL} := \Omega(A_{TL}, B_{TL}, C_{TL}) & X_{TR} := \Omega(A_{TL}, B_{BR}, C_{TR} - X_{TL}B_{TR}) \\ \hline X_{BL} := \Omega(A_{BR}, B_{TL}, C_{BL} - A_{BL}X_{TL}) & X_{BR} := \Omega(A_{BR}, B_{BR}, \\ & C_{BR} - X_{BL}B_{TR} - A_{BL}X_{TR}) \end{array} \right)$$



Partition $B \rightarrow \left(\begin{array}{c|c|c} B_{TL} & * & \\ \hline B_{BL} & B_{BR} & \end{array} \right)$, $L \rightarrow \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right)$, $G \rightarrow \left(\begin{array}{c|c} G_{TL} & 0 \\ \hline G_{BL} & G_{BR} \end{array} \right)$
 where B_{TL} , L_{TL} , and G_{TL} are 0×0

while $\text{size}(B_{TL}) < \text{size}(B)$ do

$$\left(\begin{array}{c|c|c} B_{TL} & * & \\ \hline B_{BL} & B_{BR} & \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} B_{00} & * & * \\ \hline B_{10} & B_{11} & * \\ \hline B_{20} & B_{21} & B_{22} \end{array} \right), \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} L_{00} & 0 & 0 \\ \hline L_{10} & L_{11} & 0 \\ \hline L_{20} & L_{21} & L_{22} \end{array} \right), \dots$$

Variant 1

$$\begin{aligned} G_{10} &:= B_{10} - L_{10}G_{00}^T && (\text{TRMM}) \\ G_{10} &:= G_{10}L_{00}^{-T} && (\text{TRSM}) \\ G_{11} &:= B_{11} - G_{10}L_{10}^T - L_{10}G_{10}^T && (\text{SYR2K}) \\ G_{11} &:= \text{gChol}(G_{11}, L_{11}) && (\text{gCHOL}) \end{aligned}$$

Variant 2

$$\begin{aligned} G_{10} &:= G_{10}L_{00}^{-T} && (\text{TRSM}) \\ G_{11} &:= B_{11} - G_{10}L_{10}^T - L_{10}G_{10}^T && (\text{SYR2K}) \\ G_{11} &:= \text{gChol}(G_{11}, L_{11}) && (\text{gCHOL}) \\ G_{21} &:= B_{21} - L_{21}G_{11}^T && (\text{TRMM}) \\ G_{21} &:= G_{21} - L_{20}G_{10}^T && (\text{GEMM}) \end{aligned}$$

Variant 3

...

Variant 4

...

$$\left(\begin{array}{c|c|c} B_{TL} & * & \\ \hline B_{BL} & B_{BR} & \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} B_{00} & * & * \\ \hline B_{10} & B_{11} & * \\ \hline B_{20} & B_{21} & B_{22} \end{array} \right), \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} L_{00} & 0 & 0 \\ \hline L_{10} & L_{11} & 0 \\ \hline L_{20} & L_{21} & L_{22} \end{array} \right), \dots$$

endwhile

Routine 1: gCHOL. Variant 1.

```
1 void gChol_blk_var1( FLA_Obj G, FLA_Obj L, int nb )
2 {
3     FLA_Obj GTL, GTR, GBL, GBR, G00, G01, G02, G10, G11, G12, G20, G21, G22;
4     FLA_Obj LTL, LTR, LBL, LBR, L00, L01, L02, L10, L11, L12, L20, L21, L22;
5
6     FLA_Part_2x2( G, &GTL, &GTR,
7                 &GBL, &GBR, 0, 0, FLA_TL );
8     [...]
9     while ( FLA_Obj_length( GTL ) < FLA_Obj_length( G ) ) {
10         FLA_Repart_2x2_to_3x3( GTL, GTR, &G00, &G01, &G02,
11                               &G10, &G11, &G12,
12                               GBL, GBR, &G20, &G21, &G22, nb, nb, FLA_BR );
13         [...]
14         FLA_Trmsx_external( FLA_RIGHT, FLA_LOWER_TRIANGULAR,
15                             FLA_TRANSPOSE, FLA_NONUNIT_DIAG,
16                             FLA_MINUS_ONE, G00, L10, FLA_ONE, G10);
17         FLA_Trsm( FLA_RIGHT, FLA_LOWER_TRIANGULAR, FLA_TRANSPOSE,
18                  FLA_NONUNIT_DIAG, FLA_ONE, L00, G10);
19         FLA_Syr2k( FLA_LOWER_TRIANGULAR, FLA_NO_TRANSPOSE,
20                  FLA_MINUS_ONE, G10, L10, FLA_ONE, G11 );
21         FLA_gChol_unb(G11, L11);
22     }
23 }
24 }
```

- BLAS(-like)
 - BLAS 3: GEMM, SYMM, SYRK, TRSM, ...
 - BLAS 2: GEMV, SYMV, GER, TRSV, ...
 - BLAS 1: AXPY, DOT, ...
- LAPACK
 - Factorizations
 - Inverses
- RECSY
 - Continuous-time Sylvester
 - Continuous-time Lyapunov

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 - BLAS 3: GEMM, SYMM, SYRK, TRSM, ...
 - BLAS 2: GEMV, SYMV, GER, TRSV, ...
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- LAPACK
 - Factorizations
 - Inverses
- RECSY
 - Continuous-time Sylvester
 - Continuous-time Lyapunov
- Derivatives of the above
 - Dv(Triangular solve)
 - Dv(Cholesky)
 - ...

- 1 Two Linear Algebra Compilers
- 2 CLAK
- 3 CL1CK
- 4 Contributions**

CLAK

- High-level matrix equations
- Decomposition onto building blocks
- Replicate reasoning of a human expert
- Search guided by knowledge
- Prototypes of code generators (Matlab, Fortran)

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- High-level matrix equations
- Decomposition onto building blocks
- Replicate reasoning of a human expert
- Search guided by knowledge
- Prototypes of code generators (Matlab, Fortran)

CL1CK

- Building blocks
- Full automation of FLAME's methodology
- Dynamically increases its knowledge-base
- Potential to derive entire libraries of kernels

Additional results

- Inference engine for dynamic deduction of knowledge/properties

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- Study: DSCs vs ADIFOR for differentiated BLAS and LAPACK ops

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- Inference engine for dynamic deduction of knowledge/properties
- Study: DSCs vs ADIFOR for differentiated BLAS and LAPACK ops
- OmicABEL in GenABEL (large speedups, state-of-the-art)

● CLAK

- Fabregat & Bientinesi. A domain-specific compiler for linear algebra operations. In *High-performance computing for computational science (VECPAR 2012)*.
- Fabregat & Bientinesi. Application-tailored linear algebra algorithms. *International journal of high-performance computing and applications (IJHPCA)*, 2013.

● CL1CK

- Fabregat & Bientinesi. Knowledge-based automatic generation of Partitioned Matrix Expressions. In *Computer Algebra in Scientific Computing (CASC 2011)*.
- Fabregat & Bientinesi. Automatic generation of Loop Invariants for matrix operations. *International conference in computational science and its applications (CASA 2011)*.

● Computational Biology

- Fabregat, Aulchenko & Bientinesi. Solving sequences of generalized least-squares problems on multi-threaded architectures. *Applied Mathematics and Computation journal (AMC)*. Accepted pending minor revision.
- Fabregat & Bientinesi. Computing petaflops over terabytes of data: The case of genome-wide association studies. *ACM Transactions on Mathematical Software (TOMS)*.

● Tensor Contractions

- Di Napoli, Fabregat, Quintana & Bientinesi. Towards an efficient use of the BLAS library for multilinear tensor contractions. *Applied Mathematics and Computation journal (AMC)*. Accepted pending minor revision.

- Integration of performance analysis techniques

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- Extend the scope of CLAK

- Integration of performance analysis techniques
- Algorithm analysis and code generation for parallel archs
- Extend the scope of CLAK
- Further explore the potential of our DSCs on AD

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- HPAC
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- AICES Students and Service team
- RWTH Computing Center, DAAD

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Deutsche
Forschungsgemeinschaft

DFG